Efficient Database Model to Represent Numerical Research Data of Material Flow Analysis.

Master Thesis - Computer Science

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# Introduction

## Problem Statement

Usually, the research data of Industrial Ecology (IE) is stored in .mat or .csv files on a local machine which imposes challenges like:

- Lack of an efficient data storing mechanism.
- Lack of an online platform to query and retrieve data publicly.
- Data sharing and re-using terms has not been discussed for such data.

## Approach

- Study on Relational Database (RDB)
  - MySQL database.
  - Implementation of web interface.
- Study on Semantic Web and its components
  - RDF Data Model.
  - Apache Jena TDB triple store.

## Background



## Relational Database (RDB)

- A collection of data sets organized and stored in relational tables.
- Each of these tables has one primary key column and shares at least one column (referred as foreign key) with another table to establish relationship.
- Data can be extracted by querying with query language like SQL.
- RDBMS is a database management system that is used to create and maintain relational databases.
- Popular RDBMS examples include Microsoft Access, SQL Server, Oracle Database, MySQL, PostgreSQL etc.



## **MySQL** Database Architecture

- **The application layer:** client-side layer which handles the connection string, authentication and most importantly security.
- The server layer: the brain of the overall architecture. Any kind of query statement is executed in this layer.
- **The storage engine layer:** offers different storage engines.

#### Indexes:

- Indexes are used to find a data entry quickly.
- Let us consider following example:

SELECT \* FROM table WHERE id=1;

- Without the indexes, the query will go through every row and column.
- The optimization here is to add an index, for example, primary key. It runs the query only against the table indexes rather than all the column data.

# Background

## **The Semantic Web**

- "The Semantic Web provides a common framework that allows data to be shared and reused across application, enterprise, and community boundaries".
- Data is structured and published following semantic web standards: RDF.

#### **RDF- The Data Model**

Resource Description Framework is a "framework for representing information in the Web".

- The basic elements of RDF are triples.
- A triple is a set of three entities that form a statement in the form of subjectpredicate-object expressions.



## **URIs, Vocabularies & Ontologies**

 Uniform Resource Identifiers (URIs) are short strings that identify resources in the web.

<http://dbpedia.org/resource/Berlin> ..(1) <http://www.w3.org/1999/02/22-rdf-syntax-ns#type> ..(2) <http://dbpedia.org/class/yago/CapitalsInEurope> ..(3)

> subject  $\rightarrow$  predicate  $\rightarrow$  object (1)  $\rightarrow$  (2)  $\rightarrow$  (3)

This triple refers that resource (1) has a relationship (2) with the resource (3).

- The "vocabularies" on the Semantic Web is used for data integration between data sets.
- The term "ontology" is for more complex, and more formal collection of terms, where "vocabulary" is used for more basic use.

# Background

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## **SPARQL – Query Language**

- SPARQL is the standardized query language for RDF, just like SQL is the standardized query language for relational databases.
- Like SQL, SPARQL also follow the same "SELECT...FROM...WHERE..." query structure.

```
:id1 foaf:name "André Schürrle"
:id1 foaf: based_near : Dortmund
:id2 foaf:name "Nils Petersen"
:id2 foaf:based_near : Freiburg
```

```
SELECT ?name
WHERE {
?x foaf:name ?name .
?x foaf:based_near : Freiburg .
}
```

The predicate has a constant value of foaf:based\_near and the object has a constant value of : *Freiburg* match to one of RDF triples. The result is *Nils Petersen*.

# Background

#### **Triple Stores**

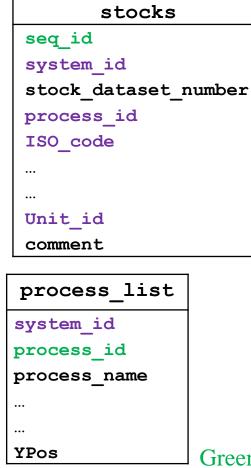
- Triple Store (RDF triple store) is a specialized DBMS for RDF triples.
- RDF data can be stored in two ways:
  - in *files,* triples are stored following one of the serialization formats, or
  - in special kind of databases for triples, called *triple stores*.
- *Triple stores* have three possible architectures:
  - In-memory: stores the triples in main memory. It is fast but expensive.
  - Native Store: storage systems with own database. For example- Jena TDB, Sesame Native, Virtuoso, AllegroGraph, Oracle 11g etc.
  - Non-native Store: storage using a third-party RDBMS. For example- Jena SDB backed by MySQL database.

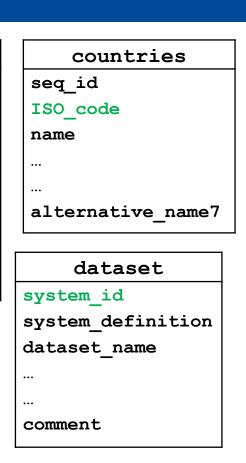
#### MySQL Design Process:

- Six different CSV files containing the data of country, unit specification, dataset (list of datasets), process list, stock data and flow data between processes.
- Database Schema, Tables, and Constraints
  - Database Schema => database name.
  - 6 tables.
  - 6 primary keys and 11 foreign keys; are said to be database constraints.
  - 589165 data rows.
  - Roughly 55mb of SQL dump.

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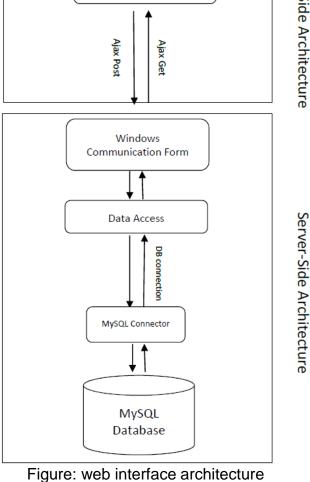
Green columns = primary keys. Purple columns = foreign keys.

flows
seq_id
system_id
flow_dataset_number
process_id_source
process_id_target
region_source
region_target
Unit_id
comment

Unit_classification
unit_id
si_unit_id
unit_code
factor

#### **Construction of Relational Database** Client-Side Architecture **Client-Side Presentation** Implementation of web interface Ajax Post Ajax Get **Back-End Development** Mechanisms to communicate with • Windows server-side database. Communication Form

- Front-End Development
  - Mechanisms to present data in clientside.



#### Data Table with the catalogue of datasets:

 Query 1: Select all the dataset available in the database from dataset table;

Select Data Set										
-	System ID 🗶	Dataset Name	Reference Date	Most Recent Update	Corresponding Auth	Other Author Info	Document Reference	Region	Time Period Of Analy	Indicator Element
	1	Global multiregional ste	05.08.2012 00:00:00	05.08.2012 00:00:00	Pauliuk, Stefan	Tao Wang, Daniel B M	10.1016/j.resconrec.20	World	1700-2008	Iron
		_				_				

Figure: DataGrid with list of datasets available in the database

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#### Data Table with material stocks:

 Query 2: Select all the stocks within the dataset (identified by the id) and filter the data with year (2007, 2008) and country (Austria, Germany, Switzerland and Italy) from stock table;

Pr	ocesses w	ith stocks	6											
Ch	Choose Stock Process Slag piles 🗸													
Filte	Filter by Year and Region													
		-												
	2008 × 2007 ×													
Tag	<b>is:</b> 2008, 2007													
	Austria × Germ	any × Switzer	land × Italy ×											
Tag	<b>is:</b> Austria, Germa	any, Switzerland,	Italy											
Sele	ct Stock Process													0
	stock dataset num	process name	country name	year	age cohort	indicator element	aspect of dataset	value	error type	error value 1	error value 2	data quality	unit id	comment
	762	Slag piles		2007		slag	Stock	14364,0962395					kt	
	763	Slag piles	Austria	2008		slag	Stock	14742,1002813					kt	
	5340	Slag piles		2007		slag	Stock	157162,243482					kt	
				2008		slag	Stock	159458,228723					kt	
			Italy	2007		slag	Stock	59069,7802223					kt	
				2008		slag		60640,8354917					kt	
	13515	Slag piles	Switzerland	2007		slag	Stock	2070,50257304					kt	
	13516	Slag piles	Switzerland	2008		slag	Stock	2136,16484003					kt	

#### Figure: Data Grid with stock data between processes

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#### Data Table with material flows:

 Query 3: Select all the flows within the dataset (identified by the id) and filter the data with year (2005) and source process 'scrap market' and target process 'foundries' from flow table;

Flows betwee	Flows between processes												
Choose Source and Tar	get Process				Scrap marke	et		~	Fo	undries			~
Filter by Year and	ilter by Year and Region												
2005 ×													
Tags: 2005													
Tags:													
Select Flow Processes													0
	rce process n target process i	a source country n	target country na	year	age cohort	indicator elemen	nt value	error type	error value 1	error value 2	data quality	unit id	comment
106 Scra	p market Foundries	Afghanistan	Afghanistan	2005		steel scrap	0					kt/yr	^
215 Scra	p market Foundries	Albania	Albania	2005		steel scrap	0					kt/yr	
324 Scra	p market Foundries	Algeria	Algeria	2005		steel scrap	69,2790266983					kt/yr	
433 Scra	p market Foundries	Angola	Angola	2005		steel scrap	10,25217303					kt/yr	
542 Scra	p market Foundries	Argentina	Argentina	2005		steel scrap	331,052613857					kt/yr	
651 Scra	p market Foundries	Australia	Australia	2005		steel scrap	473,021242061					kt/yr	
760 Scra	p market Foundries	Austria	Austria	2005		steel scrap	411,972432819					kt/yr	
869 Scra	p market Foundries	Bahrain	Bahrain	2005		steel scrap	0					kt/yr	
978 Scra	p market Foundries	Bangladesh	Bangladesh	2005		steel scrap	530,380352549					kt/yr	
1087 Scra	p market Foundries	Barbados	Barbados	2005		steel scrap	4,61843659142					kt/yr	
1196 Scra	p market Foundries	Belgium-Luxembou	Belgium-Luxembou	2005		steel scrap	532,4552064					kt/yr	
1305 Scra	p market Foundries	Bermuda	Bermuda	2005		steel scrap	0,952782763078					kt/yr	
1414 Scra	p market Foundries	Bolivia	Bolivia	2005		steel scrap	18,4313739029					kt/yr	
	n market Foundries	Brazil	Brazil	2005		steel scrap	1992 28516852					kt/vr	~
🔎 🧔 Export Selected	Rows To CSV				IN SO Pa	ge 1 of 8 ⊳>	▶1 20 ✓						View 1 - 20 of 146

Figure: Data Grid with flow data between processes

- The majority of data on the current Web is stored in relational databases.
- Semantic web is useful, especially if data from different sources has to be shared or integrated.
- Therefore, it is important to introduce mapping technologies between relational database and RDF.

D2RQ: A mapping mechanism to use RDB as RDF graphs (in addition, it exposes RDB to RDF triples)

#### **D2RQ** architecture

- A D2RQ Engine accesses a Non-RDF database.
- The mapping language describes the relationship between an ontology (or vocabulary) and a relational data model.
- RDF APIs can be embedded with Java applications via Jena or Sesame APIs.
- The RDF dump file can be stored in any triple store.
- D2RQ server provides an HTML view to explore the mapped database.

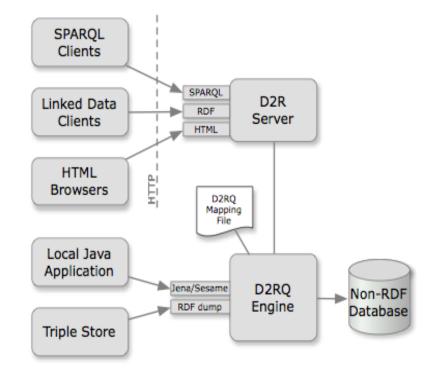


Figure: Architecture of D2RQ

#### **D2RQ: Overview and Features**

- The D2RQ platform serves different tools to offer RDF-based access to the content of relational databases.
  - The *generate-mapping* tool creates a D2RQ mapping file by analysing the relational database schema.

>generate-mapping -u username -p password -o mapping.ttl jdbc:mysql:///ief

- The *dump-rdf* tool is used to dump the contents of the entire database into a single RDF file.

>dump-rdf -f TURTLE -b http://localhost:2020/mapping.ttl > iefdatabase.ttl

- The **D2R-server** takes the mapping file as input and provides a web interface where RDF data can be browsed.

>d2r-server mapping.ttl

## Database table to RDF triples with D2RQ:

#### <"primary key column value" "tableName\_columnName" "columnValue">

id	country			process	system_ id	unit_id	value	year
17011	Germany	Stock	slag	Scrap market	1	kilo ton	444.090 7	1906



<"stocks17011"	"stocks_country" "Germany">
<"stocks17011"	"stocks_aspect_of_dataset" "Stock">
<"stocks17011"	"stocks_indicator_element" "slag">
<"stocks17011"	"stocks_process" "Scrap market">
<"stocks17011"	"stocks_system_id" "1">
<"stocks17011"	"stocks_unit_id" "kilo ton">
<"stocks17011"	"stocks_value" "444.0907">
<"stocks17011"	"stocks_year" "1906">

#### Apache Jena TDB

- TDB is a component of Jena that used as native RDF storage.
- TDB is stored in a single directory in the filing system backed by a dataset.
- A complete TDB dataset consists of:
  - **The node table:** stores the representation of RDF terms. It consists of two mappings: Node to Nodeld and Nodeld to Node.
  - Triple and Quad indexes: used for the default graph. Triples are held as 3-tuples of Nodelds in triple indexes where quads are held as 4tuples of Nodelds.
  - The prefixes table: stores index for Graph->Prefix->URI mapping. It provides the mechanism for Jena API to serialize the triples in RDF/XML or Turtle.

#### Experimental Settings:

- Two training datasets:
  - **IEF dataset:** consisting the material stock and flow data of Industrial Ecology Freiburg research group.
    - 6 tables, 589165 rows.
    - D2RQ generates roughly 10 million triples.
    - 7596007 final triples (ignoring the empty " " column values).
  - ISWC dataset: sample dataset with the information about conferences, papers, authors and topics from ISWC 2002 conference<sup>1</sup>.
    - 9 tables, 96 rows.
    - 322 triples after mapping.

<sup>1</sup>http://iswc2002.semanticweb.org/

#### **Results Evaluation:**

SQL Q	Queries on MySQL	SPARQL Queries on Jena TDB					
	Query Execution Time (ms) (average of 10 run)		Query Execution Time (ms) (average of 10 run)				
Query 1	9	Query 5	145				
Query 2	127	Query 6	1673				
Query 3	80	Query 7	11448				
Query 4	16	Query 8	10257				

#### Table 1: Query execution time comparisons for ief dataset

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#### **Results Evaluation:**

SQL Q	Queries on MySQL	SPARQL	Queries on Jena TDB
	Query Execution Time (ms) (average of 10 run)		Query Execution Time (ms) (average of 10 run)
Query 9	3	Query 13	132
Query 10	3	Query 14	7
Query 11	2	Query 15	5
Query 12	2	Query 16	12

Table 2: Query execution time comparisons for iswc dataset

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#### **Results Evaluation:**

- MySQL works faster than TDB
  - MySQL storage engines are faster architecturally.
  - Use of indexing to identify records.
- TDB works slower because of enumeration
  - TDB query execution method does not finish until results are fully enumerated.
  - Enumeration process takes longer time depending on number of predicates (number of columns in case of MySQL) and triples.

ResultSet results = qexec.execSelect(); long numResults = ResultSetFormatter.consume(results);

The variation of query execution time in TDB (when triples are available in main memory!).

## Discussions

- Most of the web contents are backed by RDB.
- RDB's are easy to implement; become complex with the number of tables increases.
- RDB works as "closed loop".
- Data can not be integrated from different sources if the database schemas are not in same structure.

- Triple Stores allow data to be extended across different data sources.
- Data can be integrated with multiple data sources by adding properties.
- A small data source can be enlarged into a bigger and richer data source.

# Discussions

## Future Work

- Query optimization for SQL joins.
- Allow filtering data with a range of years (1901-2000 instead of using 1901,1902, ...., 2000).
- Introduce filtering based on a collection of regions (not only country wise, also continents wise like "Europe").
- Improvement of default RDF vocabulary terms within D2RQ (use of more meaningful and publicly well-known vocabulary).
- Build a domain-specific complete ontology from the dataset and make it public.



# Thank You!