

# Delay-Robustness of Transfer Patterns in Public Transportation Route Planning

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Transportation Modelling, Optimization and Systems

Transfer pattern of a route

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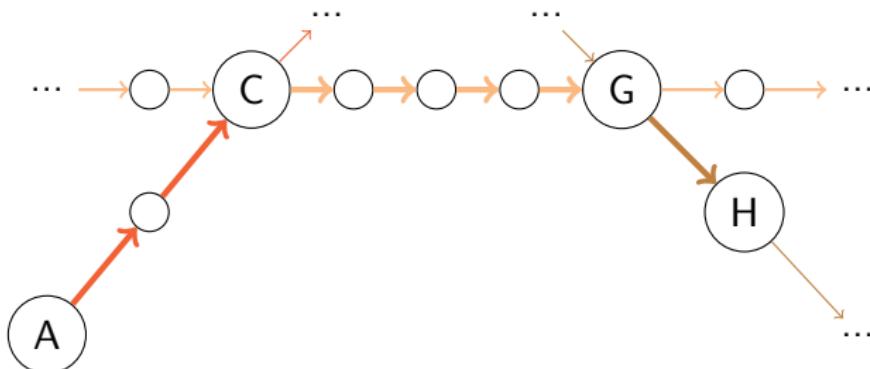
Sequence of stations where vehicle is changed

Transfer pattern of a route

:=

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Example: Route A → H has transfer pattern ACGH.



# Transfer Pattern Routing

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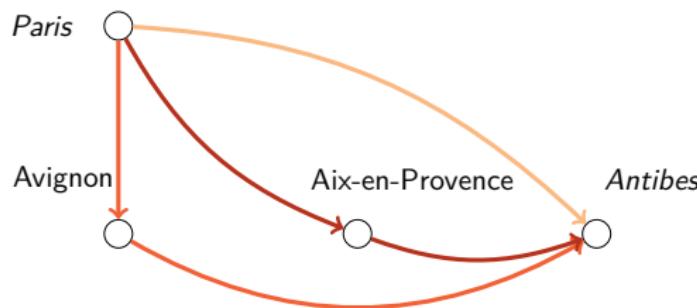


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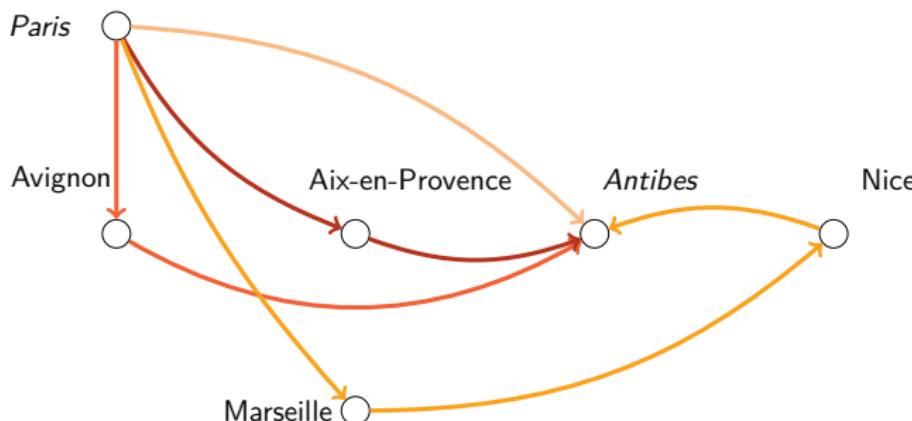


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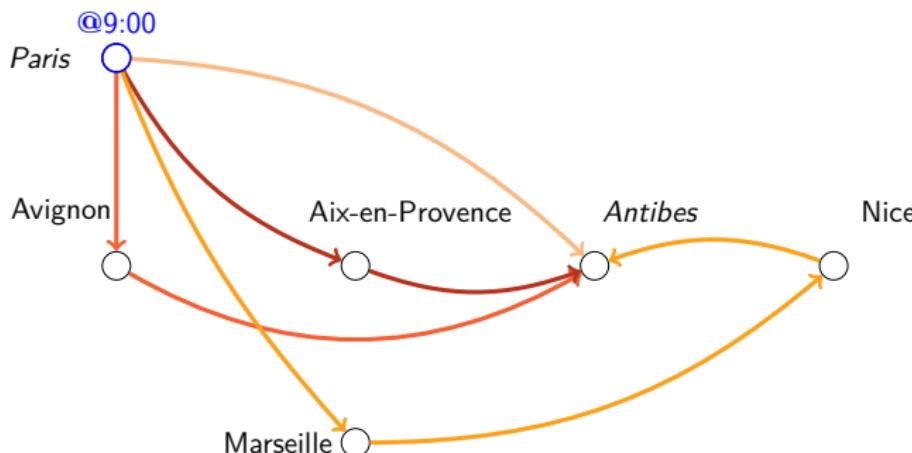


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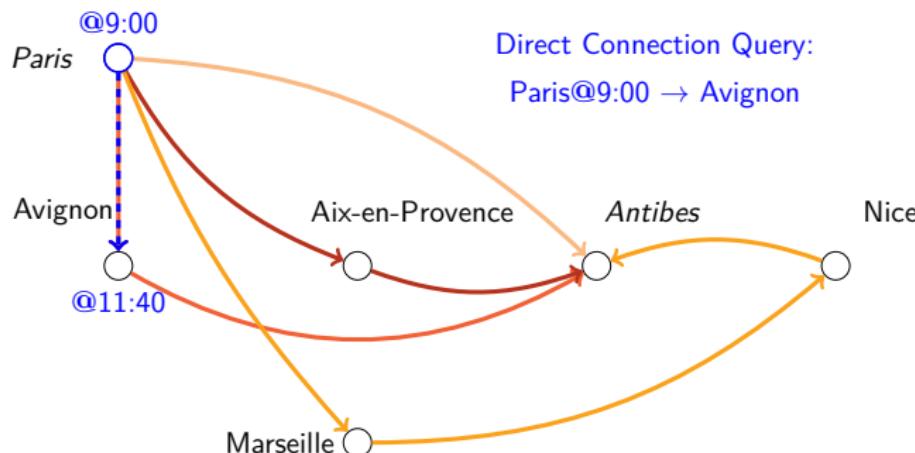


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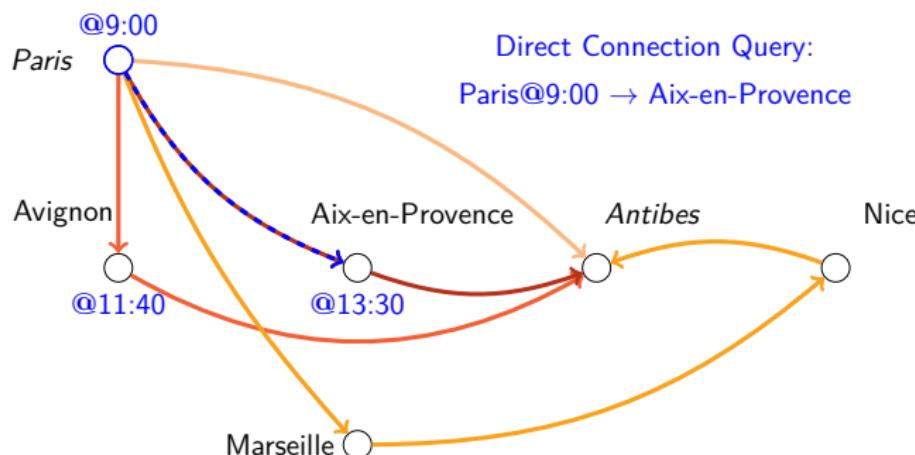


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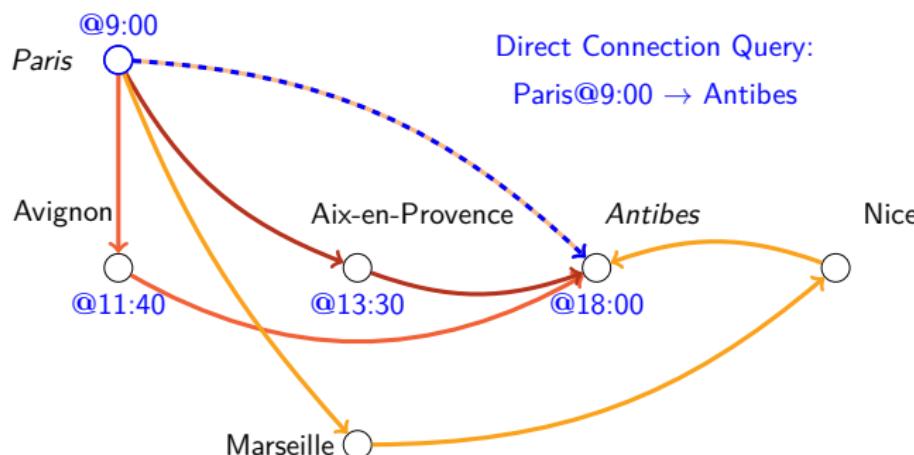


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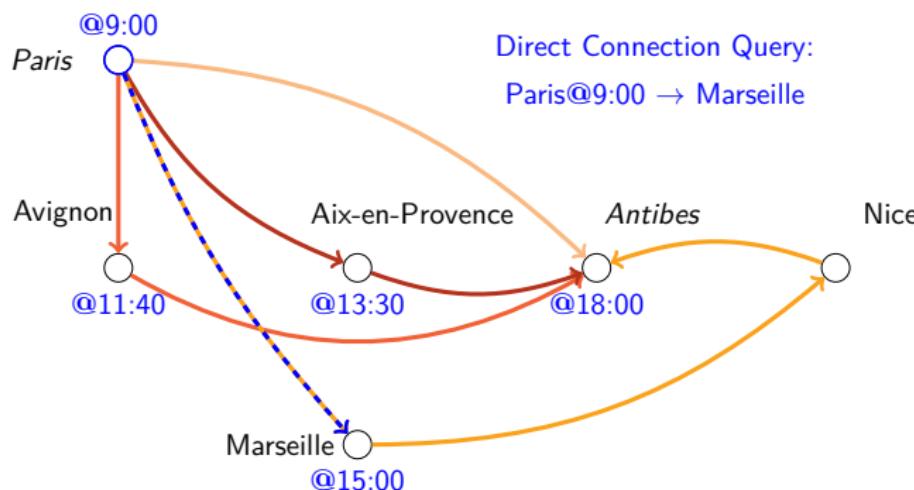


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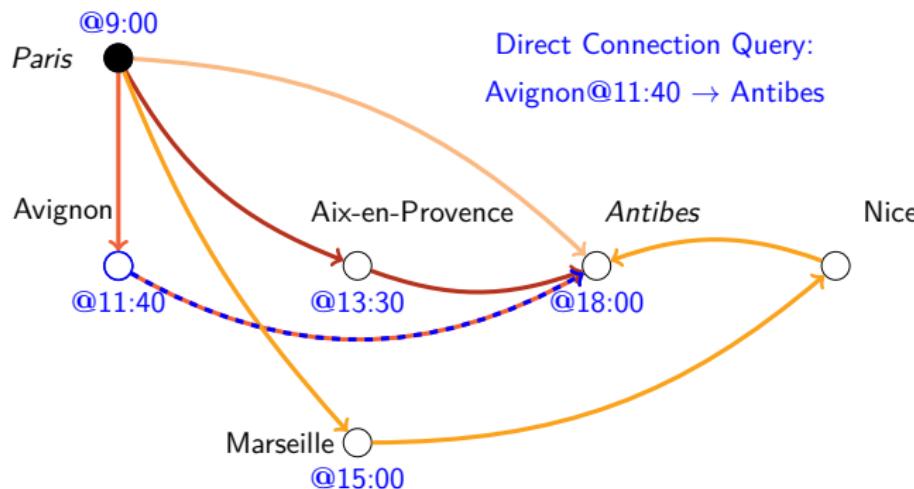


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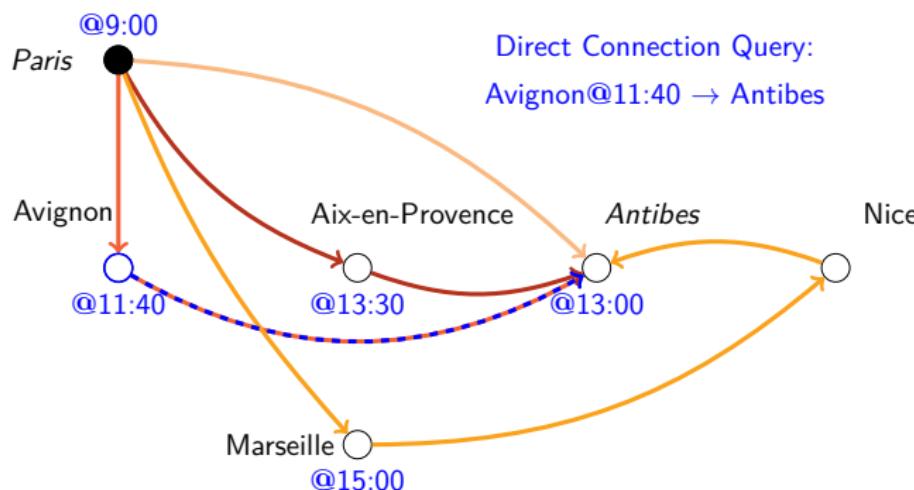


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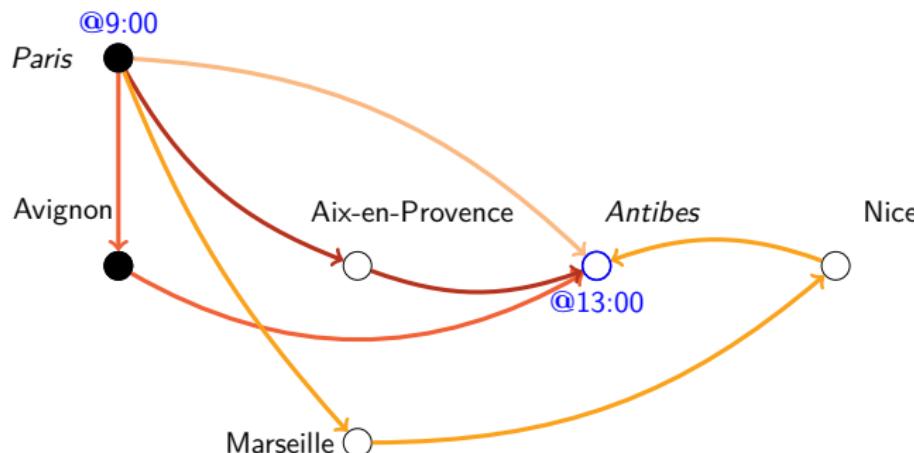


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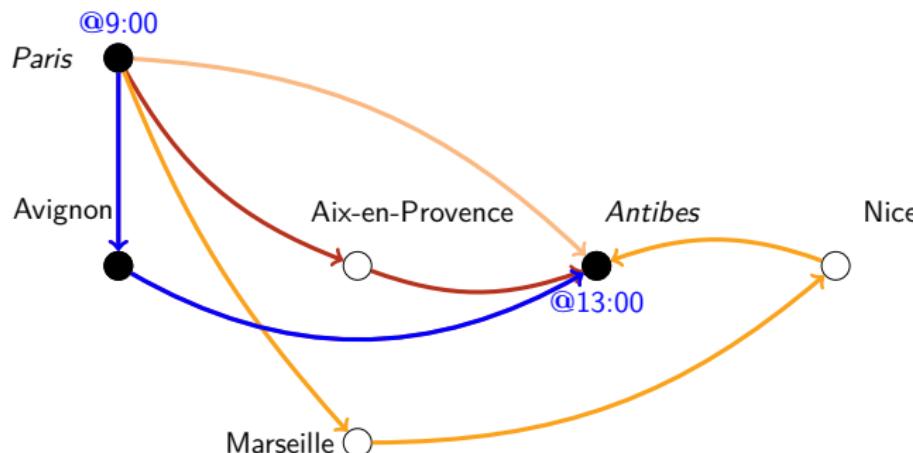


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# Transfer Pattern Routing

Public Transportation Route Planning



- ▶ Find a set of Pareto-optimal routes for a query  $A@t \rightarrow B$  given a set of timetables  $tt$
- ▶ Different structure than Road Networks, different algorithms

Transfer Pattern Routing

- ▶ *Bast et al. ESA 2010*
- ▶ Given the optimal transfer patterns, search the best route among them
- ▶ Multi-criteria
- ▶ Very fast even on huge networks (e.g. Northern America)
- ▶ State-of-the-art, Google Maps

# Transfer Pattern Routing

## Components

Data computed from timetables  $tt$

- ▶ Transfer patterns  $TP_{tt}$
- ▶ Direct connection datastructure  $DC_{tt}$

At query time

- ▶ Dijkstra on query graph built from patterns
- ▶ Relax arcs using direct connection queries



# Transfer Patterns vs. Real-Time Updates

Time-consuming precomputation

- ▶ Computation of patterns in  $O(N^2)$
- ▶ Heuristics: important stations, limits
- ▶ Still very long

Route planning applications

- ▶ Steady flow of updates during operation
- ▶ Example: Delay of trips
  - ▶ Traffic jam, strike, blocked rails, ...

No guarantee for optimal responses when trips are delayed

- ▶ Transfer patterns are computed on the original network
- ▶ Any update may introduce a new optimal pattern



# Handling real-time updates with transfer patterns

## Observation

- ▶ Typical updates introduce only small changes to the timetables,  
 $tt \rightarrow tt'$
- ▶ Computing  $TP_{tt'}$  is expensive,  $DC_{tt'}$  not.

## Main idea

- ▶ Update only the direct connection data  $DC_{tt'}$
- ▶ Search on the original transfer patterns  $TP_{tt}$

## Expectation

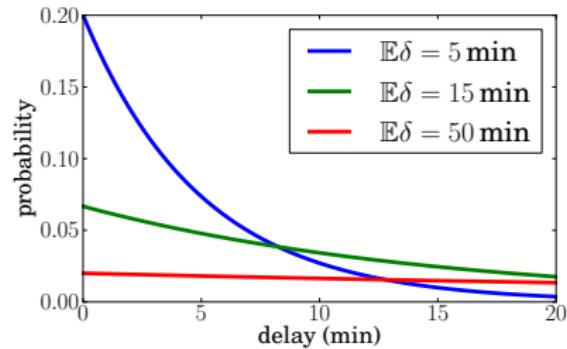
- ▶ Correct responses in most cases
- ▶ What happens elsewhere?

# How much does the quality of results suffer?

## Experimental Setup

- ▶ Compute transfer patterns
- ▶ Apply delay scenario to dataset ( $tt \rightarrow tt'$ )
- ▶ Answer N queries using  $TP_{tt}$ ,  $DC_{tt'}$
- ▶ Compare results with ground truth

Delay Scenario	Average delay $\mathbb{E}\delta$		
	5 min	15 min	50 min
LOW	25%	-	-
MEDIUM	-	25%	-
HIGH	-	-	25%
MIX LOW	10%	3%	1%
MIX NORMAL	20%	10%	5%
MIX CHAOS	40%	40%	20%



# Experiment 1

## Classification of responses

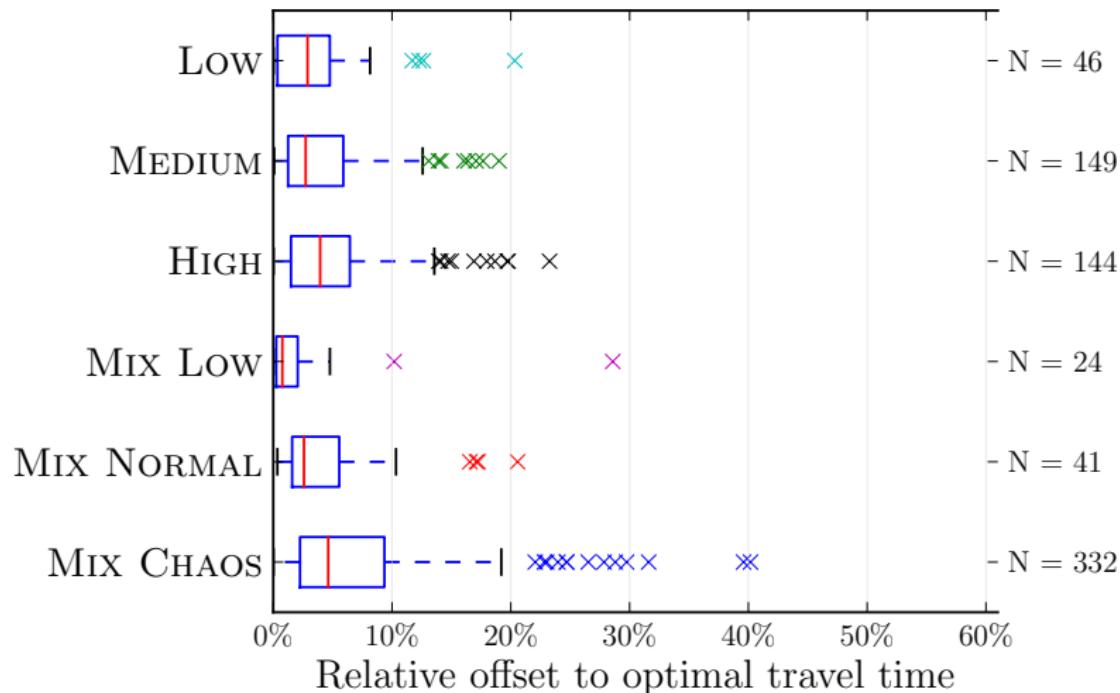
Example: 50,000 queries on New York City (2.3 M departures)

	OPTIMAL $d = 0$	ALMOST A $d \leq 5 \text{ min}$ $\wedge \frac{d}{c^*} \leq 0.05$	ALMOST B $d \leq 10 \text{ min}$ $\wedge \frac{d}{c^*} \leq 0.10$	BAD
NULL	99.99%	0.01%	0.00%	0.00%
Low	99.87%	0.09%	0.00%	0.04%
MEDIUM	99.68%	0.19%	0.03%	0.10%
HIGH	99.72%	0.17%	0.02%	0.09%
MIX LOW	99.93%	0.05%	0.00%	0.02%
MIX NORMAL	99.89%	0.07%	0.01%	0.03%
MIX CHAOS	99.19%	0.57%	0.07%	0.17%

$c := \text{travel time}$ ,  $d := c^* - c$

# Experiment 1

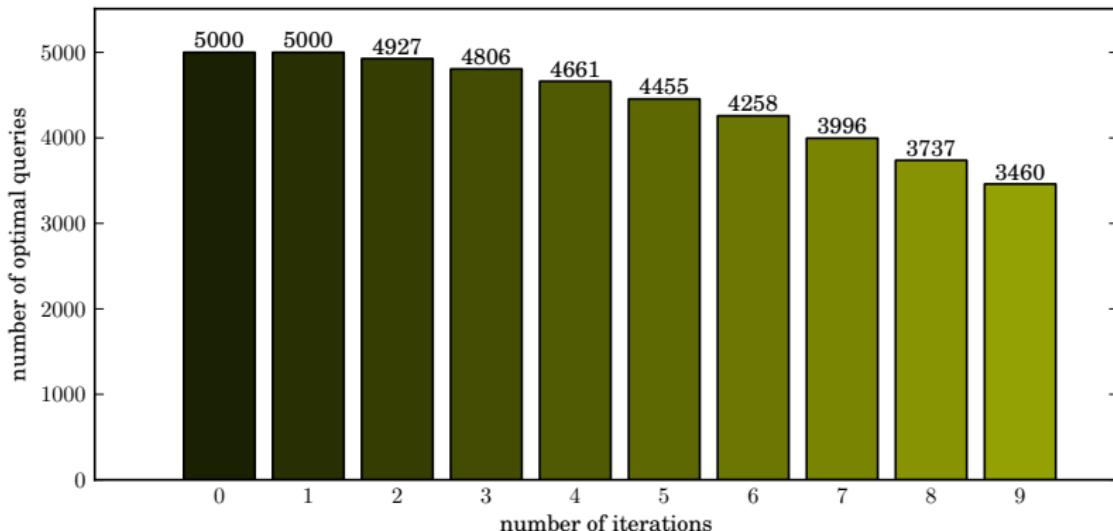
## Suboptimal responses



# Experiment 2

## Directed Delay

- ▶ Repeat a fixed query
- ▶ Delay the vehicles of the optimal connections
- ▶ Until the response becomes suboptimal



# Summary

## Delay-Robustness of Transfer Patterns

### Contribution

- ▶ Strategy to adapt to timetable updates
- ▶ Empirical proof for robustness to delay

### In the paper

- ▶ Updating the direct connection data in real-time
- ▶ Detailed experimental evaluation
- ▶ Analysis of reasons for robustness
- ▶ Discussion of possible improvements

Find it here: <http://drops.dagstuhl.de/opus/volltexte/2013/4243/pdf/5.pdf>